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Maximizing the Sum Rate in Cellular Networks Using Multiconvex Optimization

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Abstract—In this paper, we propose a novel algorithm to maximize the sum-rate in interference-limited scenarios where each user decodes its own message with the presence of unknown interferences and noise. The problem of adapting the transmit and receive filters of the users to maximize the sum-rate with a transmit power constraint is nonconvex. Our novel approach is to formulate the sum-rate maximization problem as an equivalent multiconvex optimization problem by adding two sets of auxiliary variables. An iterative algorithm, which alternatingly adjusts the system variables and the auxiliary variables is proposed to solve the multiconvex optimization problem and we show that the algorithm converges to a stationary point. The proposed algorithm is applied to a downlink cellular scenario consisting of several cells each of which contains a base station serving several mobile stations. We examine the two cases, with or without several half-duplex amplify-and-forward relays assisting the transmission. A sum power constraint at the base stations and at the relays are assumed. The applicability of our approach to the individual power constraints case is also shown. Finally, we show that the proposed multiconvex formulation of the sum-rate maximization problem is applicable to many other wireless systems in which the estimated data symbols are multiaffine functions of the system variables.

Index Terms—Sum rate maximization, interference, multiconvex function, amplify-and-forward relay.

I. INTRODUCTION

Several sophisticated solutions have been studied for future cellular systems aiming at improving both the uplink and downlink data rates. For instance, introducing multiple antennas at both base stations (BSs) and mobile stations (MSs) greatly increases the achievable rates [1]–[3]. Furthermore, employing relays in these systems extends the coverage and enhances the performance [4], [5]. However, interference is still the main performance limiting factor in cellular systems. The transmission rate, especially when the MSs are located at the cell edges, is greatly influenced by the intercell interferences. For instance for a cell edge MS, the received interference signal in the downlink can be severe and even of a comparable strength as the useful signal, which degrades the achieved rate significantly. To enhance the performance in cellular systems, smart spatial signal processing techniques at the BSs and the MSs, and also at the relays if they are employed in the system, need to be found. Apart from joint processing techniques which require data exchange among the cooperating parties, we focus on distributed signal processing techniques which require only the exchange of channel state information. If relays are exploited for interference mitigation, they can apply the distributed zero forcing technique which fully cancels the interference in the network [4]–[7]. However, this technique requires a large number of relays and relay antennas. Furthermore, cooperative relaying techniques known usually as distributed beamforming were studied extensively in the last few years [8]–[11]. Using the distributed beamforming techniques, higher performances can be achieved by increasing the number of relays. As compared to the distributed beamforming techniques which only adapt the signal processing at the relays, cooperation among the relays, BSs and MSs yields more possibilities of interference reduction and it is the main focus in this paper.

Before discussing the sum rate, we first briefly review two interference reduction techniques which have been studied extensively, i.e., interference alignment (IA) and sum mean square error (MSE) minimization. IA is achieved by aligning all the interferences in a smaller subspace of the received signal space while keeping the useful signal subspace interference free [12], [13]. IA has received great attention in the last few years [14]–[18]. Basically, the IA problem has the nice property that it is a multi-affine problem. A function is called multi-affine if there exists a partitioning of the set of variables into disjoint non-empty subsets such that the function is affine for each of these subsets of variables. Similarly, multi-convex and multi-concave functions can be defined by replacing the property of being affine by convex and concave, respectively [19], [20]. Because of the multi-affinity of the IA problem, it can be tackled by alternatively solving several linear subproblems [14], [16]. For instance, the IA problem is a multi-affine problem if relays are employed. Firstly, the filters of the BSs are optimized with fixed relay processing matrices and fixed filters at the MSs. Secondly, the relay processing matrices are

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optimized with fixed filters at the BSs and MSs. At the third step, the filters of the MSs are optimized with fixed filters at the BSs and fixed relay processing matrices. However, since IA ignores the received noise, it performs poorly at low and moderate signal to noise ratios (SNRs) [21]. On the other hand, optimizing the spatial filters at the BSs and the MSs, as well as the processing matrices at the relays if they are employed, for minimizing the sum MSE always achieves a compromise between interference reduction and noise reduction. It can be noted that only in some special cases, the sum MSE function has a convex structure [22]. In general, the sum MSE is not a convex function. However, it is a convex function of either the filters at the BS, the filters at the MSs or the relay processing matrices alone. This multi-convex structure of the sum MSE function also allows alternating minimization algorithms to achieve a stationary point [23]–[26]. Furthermore, it is worth to mention here that minimizing the sum bit error rate (BER) is an alternative objective to minimizing the sum MSE [27].

Because IA does not take noise reduction into account, it performs well only in the interference dominant regime. Furthermore, minimizing the sum MSE does not necessarily imply achieving the maximum sum rate as they are different objectives. If one considers maximizing the aggregate bit rates at the end users as a target, then sum rate maximization is a reasonable objective. Therefore, directly maximizing the sum rate is a promising goal for efficiently utilizing the limited available system resources [28]. If the interference is treated as noise and some power constraints are considered, the sum rate maximization problem is a non-convex optimization problem [29], [30]. This non-convexity of the sum rate maximization problem holds even if we optimize over either the filters at the BSs only, the filters at the MSs only or the relay processing matrices only. Therefore, iterative alternating optimization algorithms cannot be directly implemented here.

In the last decade, a lot of progress has been made in finding efficient sum rate maximization algorithms. Algorithms from global optimization theory are proposed for finding the global maximum of the sum rate [31]–[33]. Nevertheless, these algorithms suffer from high computational complexity which limits their practicality to small scenarios only. Unlike the computationally expensive global optimization algorithms, relatively low complexity suboptimum algorithms have also received great attention. Basically, the special structure of the sum rate function can be exploited to achieve a near optimum sum rate. In [34], an interference broadcast channel is considered. Instead of maximizing the sum rate, the authors maximize the product of the SNRs at the MSs. Rather than optimizing the filters at the BS and the MSs all together, it is shown that the problem can be simplified to three subproblems, which are not necessarily convex. Each subproblem aims at optimizing either the transmit powers, the BS filters or the MS filters. Geometric programming is employed for approximating the solution of the non-convex subproblems. In [35], some auxiliary variables are used to simplify the sum rate maximization problem in a broadcast channel. The authors introduce new variables to the problem such that the multiple constraints can be equivalently written as a single constraint. The sum rate function can be written as the sum of a concave function and a convex function [31], [36]. Accordingly, the authors of [36] linearly relax the convex term and solve the resulting problem iteratively. In [37], we consider a simple scenario of single antenna nodes and relays and propose a simplified formulation of the sum rate maximization problem.

Some authors also exploit the minimized MSE to maximize the sum rate. From the information theory perspective, Guo et al. have found that there is a linear relationship between the derivative of the mutual information and the minimum MSE for Gaussian channels [38]. Moreover, it is shown in [39] that this relationship holds for any wireless system with linear filters. Considering a broadcast channel scenario, the relationship between the derivative of the mutual information and the minimum MSE can be exploited by designing the receive filters such that the MSE at the receivers is minimized. In this case, the MSE will be a function of the transmit filters [40]. Accordingly, the sum rate maximization problem for optimizing the transmit filters can be formulated as a minimization of the sum of log-MSEs. An approximate solution of this new formulation is found using geometric programming [41]. Designing the receive filters to minimize the MSE and optimizing the remaining variables to maximize the sum rate is also considered in [42]. It is shown that by relaxing the sum rate maximization problem and adding some auxiliary variables, a successive convex approximation approach can be applied [42]. The authors show that the successive convex approximation algorithm converges always to a KKT point. For a broadcast channel scenario, the receive filters are designed aiming at minimizing the MSE and by adding some auxiliary variables, the sum rate maximization problem is reformulated as a biconvex optimization problem of the transmit filter and the added auxiliary variables [40]. Note that a biconvex function is a multi-convex function with two partitions [19]. This work is extended to many different scenarios such as MIMO interference channels [43], interfering broadcast channels [44], and relay interference channels [45].

In the present paper, we aim at formulating the sum rate maximization problem as a multi-convex problem so that it can be efficiently solved by iterative algorithms. We specifically consider the downlink transmission in a cellular scenario with BSs serving multiple MSs, although the same approach can be applied to the uplink transmission as well. The transmission from the BSs to the MSs takes place either through several non-regenerative relays or directly without relays. First, we focus on describing our approach for a two-hop transmission scheme where relays are employed. Then, we show that the approach can also be applied to other wireless systems by taking the single-hop transmission scheme without relays as an example. The key idea of our approach is to replace the signal to interference plus noise ratio (SINR) at a MS by a new term whose maximal value is found to be 1+SINR. Using this new term, we formulate a multi-concave objective function. We will show that this objective function has the same structure as the sum rate function and, therefore, maximizing this objective function is equivalent to maximizing the sum rate function. The weighted MSE approach [40], [43]–[45] uses the same objective as our approach, which is the sum rate maximization. However, we propose a new formulation different from the weighted MSE.
Our new formulation is equivalent to the sum rate formulation for all possible values of the system variables if the auxiliary variables are optimized whereas the weighted MSE formulation is equivalent to the sum rate function only at the points where the receive filters are designed as minimum MSE filters. These two different formulations lead to different problem decompositions. From the numerical results, we observed that our proposed algorithm performs better as compared to the weighted MSE algorithm at moderate and high SNRs as will be shown in Section VI.

The rest of this paper is organized as follows. In the next section, a two-hop transmission scenario and a single-hop transmission scenario are described. In Section III, the two-hop transmission is first investigated and based on it, the multi-convex formulation of the sum rate is derived. An iterative sum rate maximization algorithm is proposed in Section III-E. To show that our idea is quite general and fits in many scenarios, we derive the multi-convex problem formulation for the single-hop transmission in Section IV. A few additional aspects are discussed in Section V and the performance of the proposed algorithm is shown in Section VI. In Section VII, the conclusions are drawn.

II. SYSTEM MODEL

A. Two-Hop Interference Broadcast Scenario

In this paper, we will consider two related scenarios, i.e., a two-hop interference broadcast scenario and a single-hop interference broadcast scenario. The former will be described here, and the latter will be described in Section II-B.

A downlink cellular scenario consisting of \( K \) cells is considered. Each cell contains a BS with \( N_B \) antennas, and \( M \) MSs with \( N_M \) antennas each. We first assume that the direct channels between the BSs and the MSs are relatively weak due to the radio environment so that they can be neglected. To enable the communication between the BSs and the MSs, \( R \) half-duplex relays with \( N_R \) antennas each are deployed in the scenario. The transmission takes place in two subsequent time slots as illustrated in Fig. 1. In the first time slot, the BSs transmit to the relays. In the second time slot, the relays retransmit a linearly processed version of what they received in the first time slot to obtain the estimated data at the MSs. The channels between the communication parties are assumed to remain constant during the transmission. Throughout this paper, we restrict our discussion to the case where each MS receives a single desired data symbol from the corresponding BS. Accordingly, each BS transmits simultaneously \( M \) complex valued data symbols with \( M \leq N_B \).

Let \( k \in \{1, \ldots, K\} \), \( m \in \{1, \ldots, KM\} \), and \( r \in \{1, \ldots, R\} \) denote the indices of the BSs, the MSs, and the relays, respectively. Then, the data symbol transmitted by the corresponding BS for the \( m \)-th MS is denoted by \( d(m) \in \mathbb{C} \) and all the data symbols transmitted by the \( k \)-th BS are denoted by the vector \( d^{(k)} \in \mathbb{C}^M \). For each BS \( k \), the transmitted data symbols are preprocessed by a linear transmit filter denoted by \( \mathbf{V}^{(k)} \in \mathbb{C}^{N_B \times M} \). The signal transmitted by the \( k \)-th BS reads

\[
\mathbf{s}^{(k)}_{B} = \mathbf{V}^{(k)} \mathbf{d}^{(k)}. \tag{1}
\]

The received signal vector at the \( r \)-th relay is

\[
\mathbf{e}^{(r)}_{R} = \sum_{k=1}^{K} \mathbf{H}^{(r,k)}_{RB} \mathbf{s}^{(k)}_{B} + \mathbf{n}^{(r)}_{R}, \tag{2}
\]

where \( \mathbf{H}^{(r,k)}_{RB} \in \mathbb{C}^{N_R \times N_B} \) denotes the channel matrix between the \( k \)-th BS and the \( r \)-th relay, and \( \mathbf{n}^{(r)}_{R} \in \mathbb{C}^{N_R \times 1} \) represents the noises at the different antennas of the relay, which are assumed to be independently identically distributed (i.i.d.) Gaussian with zero mean and variance \( \sigma^2 \). It is assumed that the number \( N_R \) of antennas at a relay is not large enough to spatially separate the received signals, i.e., \( N_R < KM \). Therefore, the amplify and forward relaying protocol is considered. The \( r \)-th relay linearly processes its received signals with the matrix \( \mathbf{G}^{(r)} \in \mathbb{C}^{N_R \times N_R} \) and the transmitted signal of the \( r \)-th relay is denoted by

\[
\mathbf{s}^{(r)}_{R} = \mathbf{G}^{(r)} \mathbf{e}^{(r)}_{R}. \tag{3}
\]

Furthermore, the received signal vector at the \( m \)-th MS is

\[
\mathbf{e}^{(m)}_{M} = \sum_{r=1}^{R} \mathbf{H}^{(m,r)}_{MR} \mathbf{s}^{(r)}_{R} + \mathbf{n}^{(m)}_{M}, \tag{4}
\]

where \( \mathbf{H}^{(m,r)}_{MR} \in \mathbb{C}^{N_M \times N_R} \) denotes the channel matrix between the \( r \)-th relay and the \( m \)-th MS, and \( \mathbf{n}^{(m)}_{M} \in \mathbb{C}^{N_M \times 1} \) represents the noises at the MS, which are also assumed to be i.i.d. Gaussian with zero mean and variance \( \sigma^2 \). Then the \( m \)-th MS can linearly post-process its received signal vector \( \mathbf{e}^{(m)}_{M} \) using a linear receive filter \( \mathbf{u}^{(m)} \in \mathbb{C}^{N_M \times 1} \) to obtain the estimated data symbol as

\[
\hat{\mathbf{d}}^{(m)} = \mathbf{u}^{(m)} \mathbf{e}^{(m)}_{M}.
\]
\[ d^{(m)} = u^{(m)H} e^{(m)} = u^{(m)H} \left( \sum_{r=1}^{R} \sum_{k=1}^{K} H_{MR}^{(m,r)} G^{(r)} H_{RB}^{(r,k)} v^{(k)} d^{(k)} + \sum_{r=1}^{R} H_{MR}^{(m,r)} G^{(r)} n_R^{(r)} + n_M^{(m)} \right), \]  
\tag{5}

where \((\cdot)^H\) denotes the Hermitian of a matrix.

Suppose the duration of each time slot is normalized to one, it is assumed that the transmitted data symbols are uncorrelated and they have the same average power

\[ E \left\{ |d^{(m)}|^2 \right\} = P_d, \]  
\tag{6}

for all \(m = 1, \ldots, KM\) where \(E \{ \cdot \} \) denotes the expectation. The sum power constraint at the BSs is given by

\[ \sum_{k=1}^{K} \text{tr} \left( E \left\{ s_B^{(k)} s_B^{(k)H} \right\} \right) \leq P_B. \]  
\tag{7}

The sum power constraint at the relays is given by

\[ \sum_{r=1}^{R} \text{tr} \left( E \left\{ s_R^{(r)} s_R^{(r)H} \right\} \right) \leq P_R. \]  
\tag{8}

**B. Single-Hop Interference Broadcast Scenario**

The second scenario we consider is similar to the one described in Section II-A, except that the direct channels between the BSs and MSs are assumed to be usable and no relays are deployed. Accordingly, the channel between the BSs and MSs is an interfering broadcast channel [44]. The MSs receive signals directly from the BSs within a single time slot as illustrated in Fig. 2. Therefore, the received signal vector at the \(m\)-th MS reads

\[ e^{(m)} = \sum_{k=1}^{K} H_{MB}^{(m,k)} s_B^{(k)} + n_M^{(m)}, \]  
\tag{9}

where \(H_{MB}^{(m,k)} \in \mathbb{C}^{N_M \times N_B}\) denotes the channel matrix between the \(m\)-th MS and the \(k\)-th BS. Similar to (5), the estimated data symbol at the \(m\)-th MS is calculated as

\[ \hat{d}^{(m)} = u^{(m)H} \left( \sum_{k=1}^{K} H_{MB}^{(m,k)} v^{(k)} d^{(k)} + n_M^{(m)} \right). \]  
\tag{10}

Furthermore, only the power constraint (7) at the BSs is relevant for the single-hop scenario. It shall be pointed out here that our approach can be applied to the general scenario where both direct and relay links are considered. In [37], we applied our approach to a scenario consisting of \(K\) source-destination pairs and \(R\) single antenna relays. Both direct and relay links are considered.
except for the $m$-th diagonal element being zero. The received interference plus noise at the $m$-th MS is given by
\[ z^{(m)} = \sum_{r=1}^{R} H_{MR}^{(m,r)} G^{(r)} H_{RB}^{(r,k)} V^{(k)} Y^{(m)} d^{(k)} + \sum_{r=1}^{R} H_{MR}^{(m,r)} G^{(r)} n_{R}^{(r)} + n_{M}^{(m)}. \] (16)

The first term and the second term of (16) represent the received intra-cell and inter-cell interference, respectively. The last two terms of (16) describe the noises received by the $m$-th MS, including the noise retransmitted by the relays. Based on this, the receive SINR at the $m$-th MS can be written as
\[ \gamma^{(m)} = \frac{P_{d} |u^{(m)} H^{(m)} q^{(m)}|^{2}}{E \left[ |u^{(m)} H^{(m)} z^{(m)}|^{2} \right]}, \] (17)
and thus, the sum rate is calculated as
\[ C(V, G, U) = \sum_{m=1}^{K M} \log_{2} \left( 1 + \gamma^{(m)} \right), \] (18)
which is a function of the tuples of variables $V$, $G$, and $U$. For the considered two-hop transmission scheme, the sum rate maximization problem for optimizing the transmit filters, the relay processing matrices and the receive filters with the sum power constraints at the BSs and at the relays can be stated as
\[ \arg \max_{(V,G,U)} \{ C(V, G, U) \} \] (19)
subject to
\[ P_{d} \sum_{k=1}^{K} \text{tr} \left( V^{(k)} V^{(k) H} \right) \leq P_{B}, \] (20)
and
\[ \sum_{r=1}^{R} \text{tr} \left( G^{(r)} P_{d} \sum_{k=1}^{K} H_{RB}^{(r,k)} V^{(k)} V^{(k) H} H_{RB}^{(r,k) H} + \sigma^{2} I_{N_{R}} \right) G^{(r) H} \leq P_{R}, \] (21)
where the constraints of (20) and (21) follow from (7) and (8), respectively. The sum power constraint of (20) at the BSs is a convex set of the transmit filters. For fixed $V^{(k)}$, $\forall k$, the matrix $P_{d} \sum_{k=1}^{K} H_{RB}^{(r,k)} V^{(k)} V^{(k) H} H_{RB}^{(r,k) H} + \sigma^{2} I_{N_{R}}$ is positive definite, and thus, (21) is a convex set of $G^{(r)}$, $\forall r$. Similarly, (21) is a convex set of $V^{(k)}$, $\forall k$ for fixed $G^{(r)}$, $\forall r$. Based on this, the sum power constraint of (21) at the relays is a biconvex set of the transmit filters and the relay processing matrices. However, the objective function – the sum rate function – is not a concave function of the system variables $V$, $G$, and $U$ because the system variables appear in both the nominator and the denominator of the SINR expression, see (17). Therefore, the optimization problem of (19)–(21) is a non-convex problem.

As individual power constraints are frequently assumed in the literature, we now also show how the problem can be formulated assuming individual power constraints per BS and per relay. The individual power constraints at the $k$-th BS and $r$-th relay are
\[ P_{d} \text{tr} \left( V^{(k)} V^{(k) H} \right) \leq P_{B}, \] (22)
and
\[ P_{d} \text{tr} \left( G^{(r)} \sum_{k=1}^{K} H_{RB}^{(r,k)} V^{(k)} V^{(k) H} H_{RB}^{(r,k) H} G^{(r) H} \right) + \sigma^{2} \text{tr} \left( G^{(r) H} G^{(r)} \right) \leq P_{R}^{(r)}, \] (23)
respectively. Similar to the sum power constraints of (20) and (21), the individual power constraint at each BS $k$ forms a convex set, see (22), and the individual power constraint at each relay $r$ forms a biconvex set, see (23). Therefore, replacing the sum power constraints of (20) and (21) by the individual power constraints of (22) and (23) does not change the convexity of the considered optimization problem and so our approach is valid for individual power constraints as well. Throughout this paper, we assume the sum power constraints at the BSs and at the relays if not mentioned otherwise. However in Section VI, we will also present numerical results for the case of individual power constraints.

### B. Signal to Interference Plus Noise Ratio

With a closer look at the structure of the sum rate function of (18), one can observe that the achieved rate at a MS is a logarithmic function of $1 + \gamma^{(m)}$. Basically, the main difficulty of handling the SINR function of (17) is that both its nominator and denominator are functions of the system variables, see (15) and (16). In order to reformulate the optimization problem of (19)–(21) as a multi-convex optimization problem, a term related to the SINR is introduced in the following proposition.

**Proposition 1:** Let $w^{(m)} \in \mathbb{C}$ be a scaling factor which scales the $m$-th transmitted data symbol $d^{(m)}$. Then, the function
\[ \eta^{(m)} \left( w^{(m)} \right) = \frac{E \left[ |u^{(m)} d^{(m)}|^{2} \right]}{E \left[ |d^{(m)} - w^{(m)} d^{(m)}|^{2} \right]}, \] (24)
has a single maximum being equal to $1 + \gamma^{(m)}$, where $\gamma^{(m)}$ is defined in (Ref{eq:SINR}).

**Proof:** Using the function
\[ g \left( V, G, u^{(m)}, w^{(m)} \right) = E \left[ |d^{(m)} - w^{(m)} d^{(m)}|^{2} \right] = P_{d} \left| u^{(m) H} q^{(m)} - u^{(m)} \right|^{2} + E \left( |u^{(m) H} z^{(m)}|^{2} \right), \] (25)
(24) can be written as
\[
\eta^{(m)}(w^{(m)}) = \frac{P_d |w^{(m)}|^2}{g(V, G, u^{(m)}, w^{(m)})}. \tag{26}
\]
Since \(\mathcal{J}^{(m)}\) described in (5) is a multi-affine function of \(V\), \(G\), and \(u^{(m)}\), the function \(g(V, G, u^{(m)}, w^{(m)})\) described in (25) is a multi-convex function of \(V\), \(G\) and \(u^{(m)}\) for a fixed \(w^{(m)}\). By calculating the general derivative of \(\eta^{(m)}\) with respect to \(w^{(m)}\) and setting the result to zero, two stationary points can be calculated as
\[
w^{(m)}_0 = 0 \tag{27}
\]
and
\[
w^{(m)}_{\text{opt}} = \frac{P_d |u^{(m)} H q^{(m)}|^2 + E \left| u^{(m)} H z^{(m)} \right|^2}{P_d q^{(m)} H u^{(m)}}. \tag{28}
\]
By substituting (27) and (28) in (26), the values of \(\eta^{(m)}\) at \(w^{(m)}_0\) and \(w^{(m)}_{\text{opt}}\), respectively, are calculated as
\[
\eta^{(m)}(w^{(m)}_0) = 0 \tag{29}
\]
and
\[
\eta^{(m)}(w^{(m)}_{\text{opt}}) = \frac{P_d |u^{(m)} H q^{(m)}|^2 + E \left| u^{(m)} H z^{(m)} \right|^2}{E \left| u^{(m)} H z^{(m)} \right|^2} = 1 + \gamma^{(m)}. \tag{30}
\]
Considering the fact that the function \(\eta^{(m)}\) described in (24) is non-negative and
\[
\lim_{|w^{(m)}| \to \infty} \eta^{(m)} = 1, \tag{31}
\]
the function \(\eta^{(m)}\) must achieve its maximum at \(w^{(m)}_{\text{opt}}\).

The nice property of \(\eta^{(m)}\) is that just its denominator is a function of the system variables \(\mathcal{V}\), \(\mathcal{G}\), and \(u^{(m)}\), whereas for \(\gamma^{(m)}\), as defined in (17), both the nominator and the denominator are functions of the system variables.

C. Problem Reformulation

In the previous section, it has been shown that the term \(\eta^{(m)}\) is equivalent to the received SINR at the \(m\)-th MS when the scaling factor \(w^{(m)}\) is optimized using (28). Let
\[
w = \left(w^{(1)}, \ldots, w^{(KM)}\right)^T \tag{32}
\]
be a vector of the scaling factors and let the elements of \(w_{\text{opt}}\) be chosen as (28). Then, the function
\[
f_{\text{2hop}}(V, G, U, w) = \sum_{m=1}^{KM} \log_2 \left( \eta^{(m)} \right) \tag{33}
\]
is equivalent to the sum rate function of (18) in the sense that both have the same local and global maxima if \(w = w_{\text{opt}}\) holds. To show the concavity of \(f_{\text{2hop}}\) with respect to the tuples \(V, G\) and \(U\), using (26), (33) can be rewritten as
\[
f_{\text{2hop}}(V, G, U, w) = \sum_{m=1}^{KM} \log_2 \left( \frac{P_d |w^{(m)}|^2}{g(V, G, u^{(m)}, w^{(m)})} \right)
- \sum_{m=1}^{KM} \log_2 \left( g(V, G, u^{(m)}, w^{(m)}) \right). \tag{34}
\]
In (34), just the second term includes the system variables. Although the function \(g(V, G, u^{(m)}, w^{(m)})\) is a multi-convex function of \(V, G\) and \(u^{(m)}\) when \(w^{(m)}\) is fixed, \(\log_2 \left( g(V, G, u^{(m)}, w^{(m)}) \right)\) is not necessarily convex [46]. Accordingly, we aim at finding a new equivalent objective function which is linear in \(g(V, G, u^{(m)}, w^{(m)})\) such that we can exploit the fact that \(g(V, G, u^{(m)}, w^{(m)})\) is a multi-convex function of the system variables.

D. Multiconvex Problem Formulation

In this section, the optimization problem of (19)–(21) is reformulated as a multi-convex optimization problem. Let
\[
t = \left(t^{(1)}, \ldots, t^{(KM)}\right)^T \tag{35}
\]
be a vector of additional scaling factors. Then, the function
\[
b_{\text{2hop}}(V, G, U, w, t) = \sum_{m=1}^{KM} \left( \log_2 \left( \frac{P_d |w^{(m)}|^2}{g(V, G, u^{(m)}, w^{(m)})} \right) \right)
+ \log_2 \left( \frac{t^{(m)}}{\ln(2)} \right) \tag{36}
\]
is obviously a concave function of \(t\). By taking the first order derivative of \(b_{\text{2hop}}\) with respect to \(t^{(m)}\) and setting the result to zero, the optimum scaling factor \(t^{(m)}_{\text{opt}}\) is calculated as
\[
t^{(m)}_{\text{opt}} = \frac{1}{g(V, G, u^{(m)}, w^{(m)})}. \tag{37}
\]
Substituting (37) in (36) yields
\[
b_{\text{2hop}}(V, G, U, w, t_{\text{opt}}) = f_{\text{2hop}}(V, G, U, w) - \frac{KM}{\ln(2)}. \tag{38}
\]
From (38), it can be concluded that the new objective function \(b_{\text{2hop}}(V, G, U, w, t)\) is equivalent to the sum rate function in the sense that they both have the same global and local maxima if the optimum scaling factors in \(w\) and \(t\) are chosen. Moreover, the function \(b_{\text{2hop}}(V, G, U, w, t)\) has a single maximum at \(w = w_{\text{opt}}\) if \(V, G, U\), and \(t\) are fixed, and the function \(b_{\text{2hop}}(V, G, U, w, t)\) is
- a concave function of \(t\) if \(V, G, U\), and \(w\) are fixed because the logarithm is a concave monotonically increasing function,
- a concave function of \(V\) if \(t, G, U\), and \(w\) are fixed because \(g(V, G, u^{(m)}, w^{(m)})\) is a convex function of \(V\),
- a concave function of \(G\) if \(t, V, U\), and \(w\) are fixed because \(g(V, G, u^{(m)}, w^{(m)})\) is a convex function of \(G\), and...
a concave function of \( U \) if \( t, V, G, \) and \( w \) are fixed because \( g(V, G, u^{(m)}, w^{(m)}) \), \( \forall m \) is a convex function of \( U \). Accordingly, the sum rate maximization problem of (19)–(21) can be equivalently formulated as a multi-convex optimization problem stated as

\[
\left( V_{\text{opt}}, G_{\text{opt}}, U_{\text{opt}}, w_{\text{opt}}, t_{\text{opt}} \right) = \arg \max_{\left( V, G, U, w, t \right)} \left\{ b_{\text{2hop}} \left( V, G, U, w, t \right) \right\} \tag{39}
\]

subject to

\[
P_d \sum_{k=1}^{K} \text{tr} \left( V^{(k)} V^{(k)H} \right) \leq P_B \tag{40}
\]

and

\[
P_d \sum_{r=1}^{R} \text{tr} \left( G^{(r)} \sum_{k=1}^{K} H_{\text{RB}}^{(r,k)} V^{(k)} V^{(k)H} H_{\text{RB}}^{(r,k)H} G^{(r)H} \right) + \sigma^2 \sum_{r=1}^{R} \text{tr} \left( G^{(r)H} G^{(r)H} \right) \leq P_r. \tag{41}
\]

This problem is a multi-convex problem of \( V, G, U, \) and \( t \).

The vectors \( w \) and \( t \) of the scaling factors can be optimized using (28) and (37), respectively. With fixed scaling factors, just the last term of (36) is relevant for optimizing the system variables and thus, the optimization problem (39)–(41) can be stated as

\[
\left( V_{\text{min}}, G_{\text{min}}, U_{\text{min}} \right) = \arg \min_{\left( V, G, U \right)} \left\{ \sum_{k=1}^{M} \sum_{m=1}^{t} \frac{t_{(m)}^{(i)}}{\ln(2)} g \left( V, G, u^{(m)}, w^{(m)} \right) \right\} \tag{42}
\]

subject to

\[
P_d \sum_{k=1}^{K} \text{tr} \left( V^{(k)} V^{(k)H} \right) \leq P_B \tag{43}
\]

and

\[
P_d \sum_{r=1}^{R} \text{tr} \left( G^{(r)} \sum_{k=1}^{K} H_{\text{RB}}^{(r,k)} V^{(k)} V^{(k)H} H_{\text{RB}}^{(r,k)H} G^{(r)H} \right) + \sigma^2 \sum_{r=1}^{R} \text{tr} \left( G^{(r)H} G^{(r)H} \right) \leq P_r, \tag{44}
\]

where \( w \) and \( t \) are fixed whereas \( V, G, \) and \( U \) are the optimization variables. As described previously in Section III-B, the function \( g(V, G, u^{(m)}, w^{(m)}) \) is a multi-convex function of \( V, G, \) and \( u^{(m)} \) for fixed \( w^{(m)} \). Moreover, the power constraints of (43) and (44) are a convex set and a biconvex set, respectively. Based on this, the optimization problem of (42)–(44) is a multi-convex problem for fixed \( w \) and \( t \). By taking the general derivative of \( g(V, G, u^{(m)}, w^{(m)}) \) with respect to \( u^{(m)} \) and setting the result to zero, the optimum receive filter is calculated as

\[
\mathbf{u}_{\text{min}}^{(m)} = \left( P_d q^{(m)} q^{(m)H} + \mathbb{E} \left[ z^{(m)} z^{(m)H} \right] \right)^{-1} \cdot P_d w^{(m)H} q^{(m)}.
\]

By substituting (15), (16) and (25) in (42), the problem (42)–(44) is a convex quadratically constrained quadratic problem for optimizing \( V \) with fixed \( G, U, \) and \( t \). Tools from quadratic optimization can be applied to find the optimum transmit filters [47]. Similarly, with fixed \( V, U, \) and \( w \), the optimization problem (42)–(44) can be solved for the tuple \( G \) of the relay processing matrices using the conventional quadratic optimization tools.

In our previous work [37], we also introduce auxiliary variables for reformulating a sum rate maximization problem, but for a simple scenario of single antenna nodes and relays. In [37], the problem is simplified by fixing part of the Tx/Rx filter coefficients. The fixed part of the the Tx/Rx filter coefficients is fixed such that the simplified sum rate maximization problem can be reformulated as a convex optimization problem of the non-fixed Tx/Rx filter coefficients and the relay coefficients for given auxiliary variables. Because only part of the variables are optimized in [37], the iterative algorithm will not solve the original sum rate maximization problem. On the contrary, the optimization problem of (42)–(44) is a multi-convex problem, i.e., a non-convex problem, of \( V, G, \) and \( U \), and it cannot be solved in a single step as in [37], but rather by alternating optimization.

### E. Iterative Algorithm

In this section, an iterative algorithm which alternately maximizes the multi-convex objective function \( b_{\text{2hop}} (V, G, U, w, t) \) by sequentially optimizing \( V, G, U, w, \) and \( t \) is described. Let \( \epsilon \) be an arbitrarily small tolerance value. Then, the proposed algorithm can be summarized as follows:

1. set arbitrary initial values for \( w^{(0)} \) and \( t^{(0)} \)
2. set feasible initial values for \( V^{(0)} \) and \( G^{(0)} \)
3. in each iteration \( i \)
   4. calculate \( U^{(i)} \) given \( w^{(i-1)}, V^{(i-1)}, \) and \( G^{(i-1)} \) \( \triangleright \) using (45)
   5. calculate \( V^{(i)} \) given \( w^{(i-1)}, t^{(i-1)}, U^{(i)} \) and \( G^{(i-1)} \) \( \triangleright \) using quadratic optimization tools [48]
   6. calculate \( G^{(i)} \) given \( w^{(i-1)}, t^{(i-1)}, U^{(i)} \) and \( V^{(i)} \) \( \triangleright \) using quadratic optimization tools [48]
   7. calculate \( t^{(i)} \) given \( w^{(i-1)}, \) \( G^{(i)} \) and \( V^{(i)} \) \( \triangleright \) using (37)
8. calculate \( w^{(i)} \) given \( V^{(i)} \) and \( G^{(i)} \) \( \triangleright \) using (28)
9. stop if \( \left| b_{\text{2hop}} (V^{(i)}, G^{(i)}, U^{(i)}, w^{(i)}, t^{(i)}) - b_{\text{2hop}} (V^{(i-1)}, G^{(i-1)}, U^{(i-1)}, w^{(i-1)}, t^{(i-1)}) \right| \leq \epsilon \)

Based on this iterative algorithm, each step finds a global maximum of a convex subproblem with respect to either \( V, \) \( G, \) \( U, \) \( w \) or \( t \). Therefore, after each iteration of the above algorithm, the value of the objective function \( b_{\text{2hop}} (V, G, U, w, t) \)
is non-decreasing. Furthermore, based on the fact that the objective function $b_{2hop}(V, G, U, w, t)$ is continuous and regular and the constraint sets of (40)–(41) are compact, our proposed algorithm converges to a stationary point according to Theorem 4.1(c) in [49].

One may notice that our approach has some similarities with the weighted MSE approach [40], [43]–[45] since both of them aim at sum rate maximization and the optimum receive filter given in (45) is a scaled version of the minimum MSE filter. However, two major differences between the two approaches shall be pointed out. Firstly, we introduce a new variable $w^{(m)}$ which can be optimized such that the function $\gamma^{(m)}(w^{(m)})$ defined in (24) can replace $1 + \gamma^{(m)}$ in the sum rate function of (18). Because of this difference in the formulation, the multi-convex problem (42)–(44) is decomposable into five convex subproblems which are different from the four subproblems resulting from the weighted MSE formulation. In other words, the domain of our objective function is larger than the domain of the weighted MSE objective function, and thus, the two algorithms will in general converge to different stationary points even if they are identically initialized. Secondly, due to the new variable $w^{(m)}$, our multi-concave function of (36) is equivalent to the sum rate function of (18) for any arbitrary values of the system variables $V$, $G$, and $U$ with optimized auxiliary variables. On the contrary, the weighted MSE approach designs the receive filters $U$ as minimum MSE filters such that 1+SNIR can be replaced by 1/MSE. Therefore, the weighted MSE objective function is equivalent to the sum rate function of (18) only when the receive filters are designed as minimum MSE filters.

IV. SINGLE-HOP TRANSMISSION SCHEME

A. Problem Formulation

In this section, we will show that the proposed multi-convex formulation of the sum rate and the iterative algorithm can be applied to the single-hop interference broadcast scenarios described in Section II-B as well.

For (10), one can observe that $d^{(m)}$ is a bi-affine function, i.e., a multi-affine function with two partitions [19], of the tuple $V$ of transmit filters and the receive filter $u^{(m)}$. Then (10) can be rewritten as

$$d^{(m)} = u^{(m)H}\left(q^{(m)}d^{(m)} + z^{(m)}\right),$$

where

$$q^{(m)} = H_{MB}^{(m,k)}v^{(k,m)}$$

is the effective useful link corresponding to the $m$-th MS including the transmit filter vector $v^{(k,m)}$, and the effective interference plus noise received at the $m$-th MS is

$$z^{(m)} = H_{MB}^{(m,k)}v^{(k,m)}d^{(k)} + \sum_{l\neq k} H_{MB}^{(m,l)}v^{(l)}d^{(l)} + n^{(m)}_M.$$  

In (48), the first term and the second term represent the received intra-cell interference and the received inter-cell interference, respectively. The noise at the $m$-th MS is described by the last term of (48).

Substituting (47) and (48) into (17), the receive SINR at the $m$-th MS can be calculated. Then, the sum rate can be calculated as

$$C(V, U) = \sum_{m=1}^{KM} \log_2 \left(1 + \gamma^{(m)}\right)$$

and the sum rate maximization problem can be formulated as

$$\left\{V_{opt}, U_{opt}\right\} = \arg\max_{\left\{V, U\right\}} \left\{C(V, U)\right\}$$

subject to

$$P_d \sum_{k=1}^{K} \text{tr} (V^{(k)}V^{(k)H}) \leq P_B.$$  

Similar to the two-hop scenario discussed in Section III-A, this problem is non-convex.

B. Multi-convex Problem Formulation

For the single-hop transmission, the function $g(V, u^{(m)}, w^{(m)})$ described in (25) can be redefined using (46). The multi-concave objective function can be written as

$$b_{1hop}(V, U, w, t) = \sum_{m=1}^{KM} \left(\log_2 \left(P_d \left|w^{(m)}\right|^2\right) + \log_2 \left(\gamma^{(m)}\right) - \frac{\gamma^{(m)}}{\ln(2)} \text{g}(V, u^{(m)}, w^{(m)})\right).$$

Based on this, the sum rate maximization problem can be formulated as a multi-convex optimization problem stated as

$$\left\{V_{opt}, U_{opt}, w_{opt}, t_{opt}\right\} = \arg\max_{\left\{V, U, w, t\right\}} \left\{b_{1hop}(V, U, w, t)\right\}$$

subject to

$$P_d \sum_{k=1}^{K} \text{tr} (V^{(k)}V^{(k)H}) \leq P_B.$$  

This problem is a multi-convex optimization problem of $V$ and $U$ if the optimum $w$ and $t$ are a priori chosen. As described in Section III-D, the optimization problem of (53)–(54) can be solved alternatingly over $V$, $U$, $w$, and $t$. The iterative algorithm can be summarized as follows:

1: set arbitrary initial values for $w^{(0)}$ and $t^{(0)}$
2: set feasible initial values for $\gamma^{(0)}$
3: in every iteration $i$
4: calculate $U^{(i)}$ given $w^{(i-1)}$ and $V^{(i-1)}$ \(\triangleright\) using (45)
5: calculate $V^{(i)}$ given $w^{(i-1)}$, $t^{(i-1)}$ and $U^{(i)}$ \(\triangleright\) using quadratic optimization tools [48]
6: calculate $t^{(i)}$ given $w^{(i-1)}$, $V^{(i)}$ and $U^{(i)}$ \(\triangleright\) using (37)
7: calculate $w^{(i)}$ given $V^{(i)}$ and $U^{(i)}$ \(\triangleright\) using (28)
8: stop if $b_{1hop}\left(\gamma^{(i)}, U^{(i)}, w^{(i)}, t^{(i)}\right) - b_{1hop}\left(\gamma^{(i-1)}, U^{(i-1)}, w^{(i-1)}, t^{(i-1)}\right) \leq \epsilon$
V. Further Discussions

The key idea of the algorithm proposed in this paper is to maximize the multi-concave objective function, i.e., the function $b_{2\text{hop}}$ of (36) or the function $b_{1\text{hop}}$ of (52), which is equivalent to the sum rate function in the sense that they have the same maxima. It can be observed from the analysis in the previous sections that the new objective function must be a multi-concave function of the system variables as long as the function $g$ described in (25) is a multi-convex function. This only requires that each estimated data symbol $d^{(m)}$ is a multi-affine function of the system variables. Therefore, we may conclude that for a system in which the estimated data symbols are multi-affine functions of the system variables, the sum rate maximization problem can be equivalently formulated as a multi-convex optimization problem. Examples of such systems include the MIMO interference networks, the uplink of MIMO cellular networks, multiuser relay networks, etc.

Concerning the computational complexity of the proposed algorithm, alternatingly optimizing the tuples of the filters and the auxiliary variables can efficiently solve the multi-convex optimization problem. Specifically, the optimum scaling factors $w$ and $t$ can be computed in closed form using (28) and (37), respectively, given the tuples of the filters $U$, $V$, and $G$ if relays are employed. Using (45), the tuple of the receive filters $U$ can also be computed in closed form. Although updating the tuple of the transmit filters $V$, as well as the tuple of the relay processing matrices $G$ if relays are employed, is a quadratic optimization problem, it can be readily solved using standard convex optimization tools. A quantitative complexity comparison with the weighted MSE minimization algorithm [44] based on simulation results will be shown in Section VI.

Furthermore, the proposed algorithm can be implemented at a central unit. If no relays are employed, the central unit shall collect the downlink estimated channel state information (CSI) from the MSs, compute $V_{\text{opt}}$ and $U_{\text{opt}}$, and feed the results back to the BSs and the MSs. If a single relay is employed, the relay can play the role of the central unit. Especially in time division duplex systems, the reciprocity of the channels can be exploited so that the channels between the relay and the MSs can be estimated at the relay. If multiple relays are employed, the CSI is required to be exchanged among the relays.

Moreover, our proposed algorithm will do the best for any scenario regardless of the placement or the number of the relays. From an implementation point of view, one can set a minimum received SNR threshold such that the relays which are far away from a certain cell will not be considered for forwarding the useful signals of this cell.

In this paper, we only considered the case where each MS receives a single desired data symbol from the corresponding BS. If more than one data symbol is desired by each MS, one can expect that the estimated data symbols at a MS are superposed by correlated noise. If this correlation of the noise is ignored and the received data symbols are decoded symbol-wise, a lower bound of the sum rate can still be obtained using the proposed algorithm. However, this is beyond the scope of this paper and would need further investigations.

VI. Numerical Results

We first compare the performances of the proposed sum rate maximization algorithm with the weighted MSE minimization algorithm [44] and the IA algorithm [14]. A cellular scenario with $K$ cells, $M = 1$ MS per cell, and no relays is considered, where the single-hop transmission scheme is applied to transmit a single data symbol from each BS to the corresponding MS. The case with $K = 3$ cells and $N_B = N_M = 3$ antennas at each BS and MS as well as the case with $K = 4$ cells and $N_B = N_M = 4$ antennas at each BS and MS are both investigated. An individual power constraint $P_k$ is assumed for each BS and the SNR is defined as $\gamma_{SNR} = \frac{P_k}{N_0}$. I.i.d. complex Gaussian channels with the average channel gain being normalized to one are assumed. Fig. 3 shows the achieved sum rates of the three considered algorithms averaged over 100 randomly generated channels. The proposed sum rate maximization algorithm and the weighted MSE minimization algorithm terminate if the absolute difference of the achieved sum rates in two consecutive iterations is less than $10^{-3}$ or if the number of iterations exceeds 100. It can be observed that both the proposed sum rate maximization algorithm and the weighted MSE minimization algorithm outperform the IA algorithm at low SNRs as expected. Furthermore, for the considered scenarios the proposed sum rate maximization algorithm is able to achieve higher sum rates on average as compared to the weighted MSE minimization algorithm, especially at high SNRs.

In the following, we will investigate the convergence and the complexity of the proposed sum rate maximization algorithm in the three-cell scenario introduced above. To make sure that the performance advantage of the proposed algorithm does not only hold for a small number of iterations, we compare the achieved sum rates of the proposed algorithm and the weighted MSE minimization algorithm at an intentionally chosen high SNR of 50 dB for up to 100,000 iterations. The results are depicted in Fig. 4. We can observe that our algorithm converges notably faster than the weighted MSE minimization algorithm in this case.

![Fig. 3. Average sum rate per time slot as a function of the SNR $\gamma_{SNR}$ in dB for scenarios with $M = 1$ and no relays.](image-url)
Fig. 4. Average sum rate versus the number of iterations at the SNR $\gamma_{SNR} = 50$ dB for a scenario with $K = 3$, $M = 1$, $N_B = N_M = 3$, and no relays.

Fig. 5. Average sum rate per time slot as a function of the pseudo SNR $\gamma_{PSNR}$ in dB for a scenario with $K = 2$, $M = N_B = 3$, $R = 4$, and $N_R = N_M = 2$.

### TABLE I

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>$n$</th>
<th>$c_{cpu}$/s</th>
<th>$C$/bits</th>
</tr>
</thead>
<tbody>
<tr>
<td>sum rate max.</td>
<td>100</td>
<td>0.33</td>
<td>32.9</td>
</tr>
<tr>
<td>weighted MSE min.</td>
<td>733</td>
<td>1.43</td>
<td></td>
</tr>
</tbody>
</table>

The computational complexity of the two algorithms is compared in Table I. For a fair and meaningful comparison, we record the sum rate achieved by the proposed algorithm after 100 iterations at the SNR of 30 dB and the CPU time it takes, which is compared with the CPU time that the weighted MSE minimization algorithm takes to achieve the same sum rate. The results clearly show that the proposed algorithm needs not only fewer iterations, but also less CPU time as compared to the weighted MSE minimization algorithm.

Next, the performance of the proposed sum rate maximization algorithm is evaluated in a cellular scenario with $K = 2$ cells, $M = 3$ MSs per cell, $N_B = 3$ antennas at each BS, $R = 4$ relays, and $N_R = N_M = 2$ antennas at each relay and MS. The two-hop transmission scheme is applied. Concerning the channel model, we employ an i.i.d. complex Gaussian channel model with the average channel gain being normalized to one. To assess the performance of our proposed algorithm, two reference schemes are considered. Firstly, an IA algorithm is considered where the tuple $V$ of the transmit filters, the tuple $U$ of the receive filters, and the tuple $G$ of the relay processing matrices are alternatingly optimized to minimize the total interference leakage in the system. The considered IA algorithm is a direct extension of the interference leakage minimization algorithm proposed in [14] to a multiuser relay scenario. The second reference scheme is the sum MSE minimization algorithm which minimizes the sum MSE by alternatingly optimizing $V$, $G$, and $U$ [23]–[26].

Firstly, the achieved sum rate is considered as a performance measure. The performance is plotted as a function of the pseudo SNR which is defined as the ratio of the sum transmit power of all the BSs and relays to the noise variance $\sigma^2$, i.e.,

$$\gamma_{PSNR} = \frac{P_B + P_R}{\sigma^2}.$$  \(55\)

Fig. 5 shows the performances of the three considered algorithms averaged over 100 channel snapshots. It can be seen from Fig. 5 that the IA algorithm performs poorly as compared to the other two algorithms at low to moderate pseudo SNRs. On the one hand, the IA algorithm does not consider noise reduction. On the other hand, the IA algorithm does not intend to improve the received powers of the useful signals when minimizing the interferences. That is to say, the IA algorithm does not maximize the received powers of the useful signals. In the pseudo SNR region shown in Fig. 5, both the sum MSE minimization and the sum rate maximization algorithms achieve superior performance as compared to the IA algorithm. However, the sum rate maximization algorithm outperforms the sum MSE minimization algorithm on average. This shows that minimizing the sum MSE does not necessarily achieve high sum rates. At high pseudo SNRs, interferences become more harmful. As IA aims at perfectly nullifying all the interferences, the sum rates achieved by the IA algorithm increase approximately linearly with the pseudo SNRs and the slope is related to the achieved degrees of freedom (DoFs). Furthermore, if the sum MSE minimization and the sum rate maximization algorithms are able to find the global optima, all three curves should have the same slope at high pseudo SNRs. However, as the total available power increases, the feasible region described by the constraint sets of (20) and (21) enlarges and this complicates the search for a good stationary point for both the sum MSE minimization and the sum rate maximization algorithms. As a result, both algorithms cannot achieve the same DoFs as the IA algorithm.

Next we will take a closer look at the convergence of the proposed sum rate maximization algorithm. In Fig. 6, the approximated probability density of the sum rates achieved by the proposed sum rate maximization algorithm, the sum MSE minimization algorithm, and the IA algorithm at a pseudo SNR of $50$ dB.
of 30 dB are shown. One can observe that the IA algorithm sometimes achieves a high sum rate, but the average performance remains low. This implies that the SNR at each MS may vary across a wide range depending on the channel realization. The performance of the sum MSE minimization algorithm is more stable than that of the IA algorithm because the received useful signal powers are forced close to the transmit signal powers, i.e., the gains of the useful links are close to one. Finally, the proposed sum rate maximization algorithm achieves the highest average sum rate with the smallest variance among the three considered algorithms in this case. For a randomly given channel realization, the algorithm converges, with high probability, to a solution which achieves a sum rate in the range between 14 bits per channel use and 17 bits per channel use. However, for some channel realizations, the algorithm may also converge to solutions achieving a sum rate of about 13 bits per channel use or 18 bits per channel use. The reason is that the sum rate maximization algorithm is not guaranteed to achieve a global maximum. In fact, alternatingly adapting the sets of optimization variables may result in that one or several users are turned off. In our simulation results for instance, it may happen that zero, one, or even two of the six MSs are turned off depending on the pseudo SNR and the channel realizations. Because of this, the IA algorithm can even outperform the proposed sum rate maximization algorithm at very high pseudo SNRs.

Fig. 7 shows the average sum rate versus the number of iterations at the pseudo SNR $\gamma_{SNR} = 30$ dB for a scenario with $K = 2$, $M = N_B = 3$, $R = 4$, and $N_R = N_M = 2$.

VII. Conclusion

In this paper, the sum rate maximization problem in cellular networks is considered. It is shown that by adding two sets of auxiliary variables, this problem can be formulated as a multi-convex optimization problem. The property of multi-convexity in the new formulation makes it possible to find a stationary point using an efficient iterative algorithm. The new proposed multi-convex formulation is not limited to our considered scenario, but it can be applied to many multiuser wireless system in which the estimated data symbols are multi-affine functions of the system variables.

REFERENCES

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